

Introduction to belief functions

Thierry Denœux

Université de technologie de Compiègne, France
Institut Universitaire de France
<https://www.hds.utc.fr/~tdenoeux>

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Contents of this lecture

- 1 Fundamental concepts: belief, plausibility, commonality, conditioning, basic combination rules.
- 2 Some more advanced concepts: informational ordering, cautious rule, compatible frames.

Theory of belief functions

History

- A formal framework for representing and reasoning with uncertain information.
- Also known as **Dempster-Shafer (DS) theory** or **Evidence theory**.
- Originates from the work of Dempster (1967)¹ in the context of **statistical inference**.
- Formalized by Shafer (1976)² as a **theory of evidence**.
- Popularized and developed by Smets in the 1980's and 1990's as the **"Transferable Belief Model"**.
- Starting from the 1990's, **growing number of applications** in information fusion, knowledge representation, machine learning (classification, clustering), reliability and risk analysis, etc.

¹A. P. Dempster. Upper and lower probabilities induced by a multivalued mapping, *Annals of Mathematical Statistics*, 38:325–339, 1967.

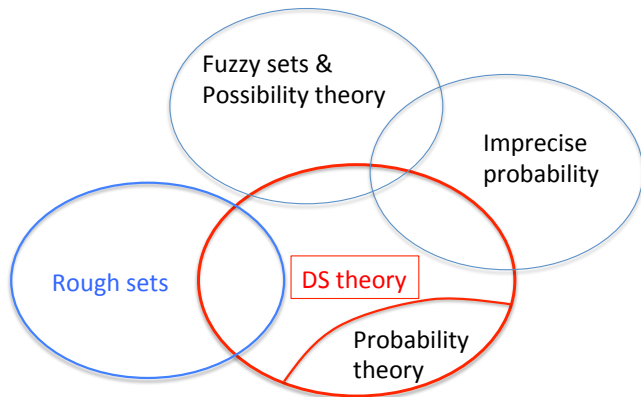
²G. Shafer. *A mathematical theory of evidence*. Princeton University Press, Princeton, N.J., 1976.

Theory of belief functions

Main idea

- The theory of belief functions extends both **logical/set-based** formalisms (such as Propositional Logic and Interval Analysis) and **Probability Theory**:
 - ▶ A belief function may be viewed both as a **generalized set** and as a **nonadditive measure**
 - ▶ The theory includes extensions of **probabilistic notions** (conditioning, marginalization) and **set-theoretic notions** (intersection, union, inclusion, etc.).
- DS reasoning produces the same results as probabilistic reasoning or interval analysis when provided with the same information.
- However, the **greater expressive power** of the theory of belief functions allows us to represent what we know in a more faithful way.

Relationships with other theories



Outline

- 1 Basic notions
 - Mass functions
 - Belief and plausibility functions
 - Dempster's rule
- 2 Selected advanced topics
 - Informational orderings
 - Cautious rule
 - Compatible frames

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Mass function

Definition

Definition (Frame of discernment, mass function, focal set)

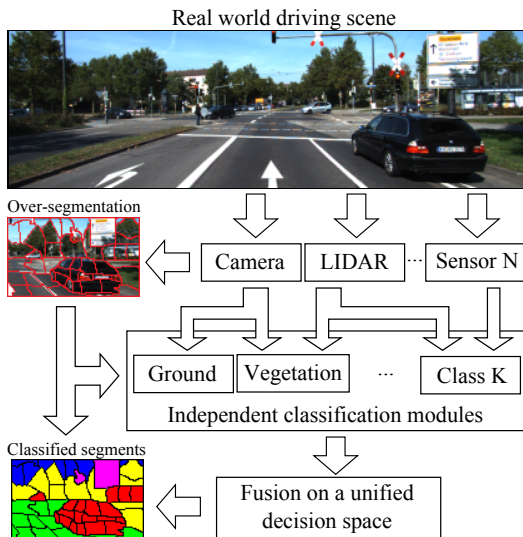
Let Ω be the finite set called a **frame of discernment**. A **mass function** on Ω is a mapping $m : 2^\Omega \rightarrow [0, 1]$ such that

$$\sum_{A \subseteq \Omega} m(A) = 1$$

Every subset A of Ω such that $m(A) > 0$ is a **focal set** of m . If $m(\emptyset) = 0$, m is said to be **normalized** (assumed in this lecture).

In DS theory, a mass function is used to represent **evidence about an uncertain variable X** taking values in Ω .

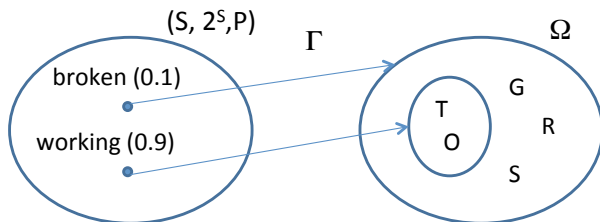
Example: road scene analysis



Example: road scene analysis (continued)

- Let X be the type of object in some region of the image, and $\Omega = \{G, R, T, O, S\}$, corresponding to the possibilities **G**rass, **R**oad, **T**ree/Bush, **O**bstacle, **S**ky.
- Assume that a lidar sensor (laser telemeter) returns the information $X \in \{T, O\}$, but we there is a probability $p = 0.1$ that the information is not reliable (because, e.g., the sensor is out of order).
- How to represent this information by a mass function?

Formalization



- Here, the probability p is not about X , but about the state of a sensor.
- Let $S = \{\text{working}, \text{broken}\}$ the set of possible sensor states.
 - ▶ If the state is “working”, we know that $X \in \{T, O\}$.
 - ▶ If the state is “broken”, we just know that $X \in \Omega$, and nothing more.
- This uncertain evidence can be represented by the following mass function m on Ω :

$$m(\{T, O\}) = 0.9, \quad m(\Omega) = 0.1$$

Meaning of a mass function

- In the previous example,
 - ▶ $m(\{T, O\}) = 0.9$ is the probability of knowing only that $X \in \{T, O\}$, and nothing more
 - ▶ $m(\Omega) = 0.1$ is the probability of knowing nothing at all.
- In general, what is the meaning (semantics) of a mass function in DS theory?
- A precise interpretation was proposed by Shafer (1981)³: random code semantics.

³G. Shafer. Constructive probability. *Synthese*, 48(1):1–60, 1981.

Random code semantics

- We consider a situation in which we receive a **coded message** containing reliable information about variable $X \in \Omega$.
- The message was encoded using some code in the set $S = \{c_1, \dots, c_n\}$.
- There is a **multi-valued mapping** $\Gamma : S \rightarrow 2^\Omega \setminus \{\emptyset\}$ that defines the meaning of the message: if code c_i was used, then the meaning of the message is “ $X \in \Gamma(c_i)$ ”.
- We don't know which code was used, but we know that each code c_i had a chance p_i of being selected, with $\sum_{i=1}^n p_i = 1$.
- Then $m(A)$ is the **probability that the meaning of the message is “ $X \in A$ ”**:

$$m(A) = P(\{c \in S : \Gamma(c) = A\}) = \sum_{i: \Gamma(c_i)=A} p_i$$

Random code semantics (continued)

- In practice, we do not receive randomly coded messages, but we can construct a mass function by **comparing our evidence about some variable X , to a hypothetical situation in which we receive a randomly coded message.**
- A mass function m can be elicited by finding the “coded-message” canonical example that is the most similar to our evidence.
- Remark: The tuple $(S, 2^S, P, \Omega, 2^\Omega, \Gamma)$ is called a **random set**. This notion plays an important role for defining belief functions in infinite spaces. I will also introduce the more general notion of **random fuzzy set** in a later lecture.

Special mass functions

Definition (Logical mass function)

If a mass function has only one focal set $A \subseteq \Omega$, it is said to be **logical**; we denote it as $m_{[A]}$. It represents “infallible” evidence that tells us that $X \in A$ for sure and nothing more. (There is a one-to-one correspondence between logical mass functions and nonempty sets).

Definition (Vacuous mass function)

The **vacuous** mass function m_γ is the logical mass function such that $m_\gamma(\Omega) = 1$. It represents **total ignorance**.

Definition (Bayesian mass function)

A mass function is **Bayesian** if its focal sets are singletons. It is equivalent to a probability distribution.

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- 1 **Basic notions**
 - Mass functions
 - **Belief and plausibility functions**
 - Dempster's rule
- 2 **Selected advanced topics**
 - Informational orderings
 - Cautious rule
 - Compatible frames

Definitions

Definition (Belief, plausibility, contour functions)

Given a mass function m on Ω , the corresponding **belief and plausibility functions** are mappings from 2^Ω to $[0,1]$ defined as follows:

$$Bel(A) = \sum_{B \subseteq A} m(B)$$

$$Pl(A) = \sum_{B \cap A \neq \emptyset} m(B) = 1 - Bel(\bar{A}).$$

The mapping $pl : 2^\Omega \rightarrow [0,1]$ such that $pl(\omega) = Pl(\{\omega\})$ is called the **contour function** associated to m .

Interpretation:

- $Bel(A)$ is a measure of the **total support** given to A
- $Pl(A)$ is a measure of the **lack of support** given to \bar{A}

Road scene analysis example

- We had $\Omega = \{G, R, T, O, S\}$ and

$$m(\{T, O\}) = 0.9, \quad m(\Omega) = 0.1$$

- Degrees of belief and plausibility of some subsets of Ω :

A	\emptyset	$\{T\}$	$\{O\}$	$\{T, O\}$	$\{T, O, R\}$	$\{T, R\}$	$\{R, S\}$	Ω
$Bel(A)$	0	0	0	0.9	0.9	0	0	1
$Pl(A)$	0	1	1	1	1	1	0.1	1

Elementary properties

- $Bel(\emptyset) = Pl(\emptyset) = 0$
- $Bel(\Omega) = Pl(\Omega) = 1$
- **Superadditivity** of Bel :

$$Bel(A \cup B) \geq Bel(A) + Bel(B) - Bel(A \cap B)$$

- **Subadditivity** of Pl :

$$Pl(A \cup B) \leq Pl(A) + Pl(B) - Pl(A \cap B)$$

- When m is Bayesian, the two mappings Bel and Pl are equal and additive:

$$Bel(A) = Pl(A) = \sum_{\omega \in A} m(\{\omega\})$$

for all $A \subseteq \Omega$.

Characterization of belief functions

- Function $Bel : 2^\Omega \rightarrow [0, 1]$ is **completely monotone**: for any $k \geq 2$ and for any family A_1, \dots, A_k in 2^Ω :

$$Bel \left(\bigcup_{i=1}^k A_i \right) \geq \sum_{\emptyset \neq I \subseteq \{1, \dots, k\}} (-1)^{|I|+1} Bel \left(\bigcap_{i \in I} A_i \right)$$

- Conversely, to any completely monotone set function Bel such $Bel(\emptyset) = 0$ and $Bel(\Omega) = 1$ corresponds a unique mass function m such that:

$$m(A) = \sum_{\emptyset \neq B \subseteq A} (-1)^{|A|-|B|} Bel(B), \quad \forall A \subseteq \Omega.$$

Relations between m , Bel and Pl

- Let m be a mass function, Bel and Pl the corresponding belief and plausibility functions.
- For all $A \subseteq \Omega$,

$$Bel(A) = 1 - Pl(\bar{A})$$

$$m(A) = \sum_{\emptyset \neq B \subseteq A} (-1)^{|A|-|B|} Bel(B)$$

$$m(A) = \sum_{B \subseteq A} (-1)^{|A|-|B|+1} Pl(\bar{B})$$

- m , Bel and Pl are thus **three equivalent representations** of a piece of evidence.

Relationship with Possibility Theory

- When the focal sets of m are nested: $A_1 \subset A_2 \subset \dots \subset A_r$, m is said to be **consonant**.
- The following relations then hold:

$$Pl(A \cup B) = \max(Pl(A), Pl(B)), \quad \forall A, B \subseteq \Omega$$

and the plausibility function can be computed from the **contour function** as

$$Pl(A) = \max_{\omega \in A} pl(\omega), \quad \forall A \subseteq \Omega$$

- Pl is then called a **possibility measure**, and Bel is the dual **necessity measure**.
- In a sense, the theory of belief functions can thus be considered as **more expressive** than possibility theory (but the combination operations are different, as we will see later).

Relation with imprecise probabilities

- A probability measure P on Ω is said to be **compatible** with m if

$$\forall A \subseteq \Omega, \quad Bel(A) \leq P(A) \leq Pl(A)$$

- The set $\mathcal{P}(m)$ of probability measures compatible with m is called the **credal set** of m

$$\mathcal{P}(m) = \{P : \forall A \subseteq \Omega, Bel(A) \leq P(A)\}$$

- Bel is the **lower envelope** of $\mathcal{P}(m)$

$$\forall A \subseteq \Omega, \quad Bel(A) = \min_{P \in \mathcal{P}(m)} P(A)$$

- Not all lower envelopes of sets of probability measures are belief functions.
- The theory of belief functions is not a theory of imprecise probabilities (the two theories have different conditioning operations).

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Road scene example continued

- Variable X was defined as the type of object in some region of the image, and the frame was $\Omega = \{G, R, T, O, S\}$, corresponding to the possibilities **G**rass, **R**oad, **T**ree/Bush, **O**bstacle, **S**ky
- A lidar sensor gave us the following mass function:

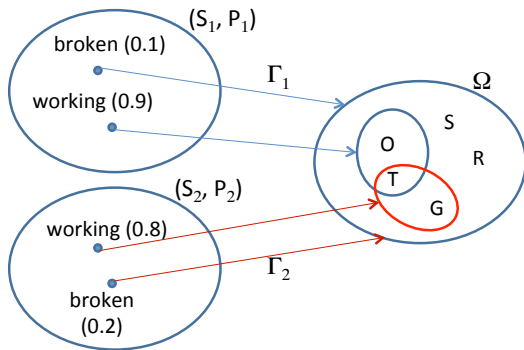
$$m_1(\{T, O\}) = 0.9, \quad m_1(\Omega) = 0.1$$

- Now, assume that a camera returns the mass function:

$$m_2(\{G, T\}) = 0.8, \quad m_2(\Omega) = 0.2$$

- How to combine these two pieces of evidence?

Analysis



- If the two sensors are in states s_1 and s_2 , then $X \in \Gamma_1(s_1) \cap \Gamma_2(s_2)$.
- If the two pieces of evidence are **independent**, then the probability that the sensors are in states s_1 and s_2 is $P_1(\{s_1\})P_2(\{s_2\})$.

Computation

$m_1 \setminus m_2$	$\{T, G\}$ (0.8)	Ω (0.2)
$\{O, T\}$ (0.9)	$\{T\}$ (0.72)	$\{O, T\}$ (0.18)
Ω (0.1)	$\{T, G\}$ (0.08)	Ω (0.02)

We then get the following combined mass function:

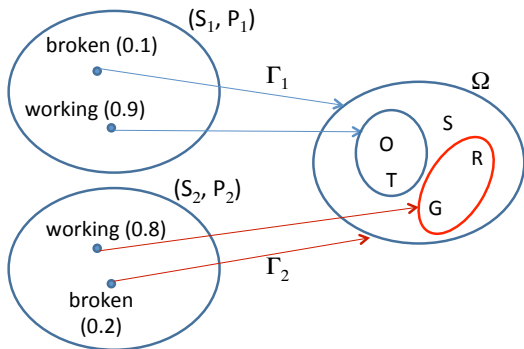
$$m(\{T\}) = 0.72$$

$$m(\{O, T\}) = 0.18$$

$$m(\{T, G\}) = 0.08$$

$$m(\Omega) = 0.02$$

Case of conflicting pieces of evidence



- If $\Gamma_1(s_1) \cap \Gamma_2(s_2) = \emptyset$, we know that the pair of states (s_1, s_2) cannot have occurred.
- The joint probability distribution on $S_1 \times S_2$ must be conditioned to eliminate such pairs.

Computation

$m_1 \setminus m_2$	$\{G, R\}$ (0.8)	Ω (0.2)
$\{O, T\}$ (0.9)	\emptyset (0.72)	$\{O, T\}$ (0.18)
Ω (0.1)	$\{G, R\}$ (0.08)	Ω (0.02)

We then get the following combined mass function,

$$\begin{aligned}
 m(\emptyset) &= 0 \\
 m(\{O, T\}) &= 0.18/0.28 = 9/14 \\
 m(\{G, R\}) &= 0.08/0.28 = 4/14 \\
 m(\Omega) &= 0.02/0.28 = 1/14
 \end{aligned}$$

Dempster's rule

Definition (Degree of conflict)

Let m_1 and m_2 be two mass functions. Their **degree of conflict** is

$$\kappa = \sum_{B \cap C = \emptyset} m_1(B)m_2(C)$$

Definition (Orthogonal sum)

Let m_1 and m_2 be two mass functions such that $\kappa < 1$. Their **orthogonal sum** is the mass function defined by

$$(m_1 \oplus m_2)(A) = \frac{\sum_{B \cap C = A} m_1(B)m_2(C)}{1 - \kappa}$$

for all $A \neq \emptyset$ and $(m_1 \oplus m_2)(\emptyset) := 0$.

Properties

Proposition

- ① If several pieces of evidence are combined, *the order does not matter*:

$$m_1 \oplus m_2 = m_2 \oplus m_1$$

$$m_1 \oplus (m_2 \oplus m_3) = (m_1 \oplus m_2) \oplus m_3$$

- ② A mass function m is *not changed if combined with the vacuous mass function $m_?$* :

$$m \oplus m_? = m$$

- ③ Let pl_1 , pl_2 and pl_{12} be the contour functions associated with, respectively, m_1 , m_2 and $m_1 \oplus m_2$. We have

$$pl_{12} = \frac{1}{1 - \kappa} pl_1 pl_2$$

Misconception about Dempster's rule

- Following a 1979 report by Zadeh, it is repeated that “Dempster's rule yields counterintuitive results” (which is usually used as a justification to introduce new combination rules)
- Zadeh's example: $\Omega = \{a, b, c\}$, two experts

$$m_1(\{a\}) = 0.99, \quad m_1(\{b\}) = 0.01 \quad m_1(\{c\}) = 0$$

$$m_2(\{a\}) = 0, \quad m_2(\{b\}) = 0.01 \quad m_2(\{c\}) = 0.99$$

We get $(m_1 \oplus m_2)(\{b\}) = 1$, which is claimed to be “counterintuitive” because both experts considered b as very unlikely.

- But Expert 1 claims that c is absolutely impossible, and Expert 2 claims that a is absolutely impossible, so b is the only remaining possibility!
- Dempster's rule does produce sound results when used and interpreted correctly.

Dempster's conditioning

- **Conditioning** is a special case of Dempster's rule, where a mass function m is combined with a logical mass function $m_{[A]}$. Notation:

$$m \oplus m_{[A]} = m(\cdot | A)$$

- It can be shown that

$$PI(B | A) = \frac{PI(A \cap B)}{PI(A)}.$$

- Generalization of **Bayes' conditioning**: if m is a Bayesian mass function and $m_{[A]}$ is a logical mass function, then $m \oplus m_{[A]}$ is a Bayesian mass function corresponding to the conditioning of m by A .

Commonality function

- **Commonality function:** let $Q : 2^\Omega \rightarrow [0, 1]$ be defined as

$$Q(A) = \sum_{B \supseteq A} m(B), \quad \forall A \subseteq \Omega$$

- Conversely,

$$m(A) = \sum_{B \supseteq A} (-1)^{|B \setminus A|} Q(B)$$

- Q is another equivalent representation of a belief function.

Commonality function and Dempster's rule

- Let Q_1 and Q_2 be the commonality functions associated to m_1 and m_2 .
- Let Q_{12} be the commonality function associated to $m_1 \oplus m_2$.
- We have

$$Q_{12}(A) = \frac{1}{1 - \kappa} Q_1(A) Q_2(A), \quad \forall A \subseteq \Omega, A \neq \emptyset$$
$$(Q_1 \oplus Q_2)(\emptyset) = 1$$

Smets' disjunctive rule

- Let m_1 and m_2 be two mass functions induced by random messages/sets $(S_1, 2^{S_1}, P_1, \Omega, 2^\Omega, \Gamma_1)$ and $(S_2, 2^{S_2}, P_2, \Omega, 2^\Omega, \Gamma_2)$.
- Previously, we have assumed that **both messages were reliable**, i.e., if the true codes are $c_1 \in S_1$ and $c_2 \in S_2$, we can conclude that $X \in \Gamma_1(c_1) \cap \Gamma_2(c_2)$ for sure.
- We can weaken this assumption by supposing only that **at least one of the two messages is reliable**, i.e., if the true codes are $c_1 \in S_1$ and $c_2 \in S_2$, we can only conclude that $X \in \Gamma_1(c_1) \cup \Gamma_2(c_2)$ for sure.
- This leads to the **Smets' disjunctive rule**:

$$(m_1 \circledast m_2)(A) = \sum_{B \cup C = A} m_1(B) m_2(C), \quad \forall A \subseteq \Omega$$

- $Bel_1 \circledast Bel_2 = Bel_1 \cdot Bel_2$

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Informational comparison of belief functions

- Let m_1 and m_2 be two mass functions on Ω
- In what sense can we say that m_1 is **more informative (committed)** than m_2 ?
- Special case:
 - ▶ Let $m_{[A]}$ and $m_{[B]}$ be two logical mass functions
 - ▶ $m_{[A]}$ is more committed than $m_{[B]}$ iff $A \subseteq B$
- Extension to arbitrary mass functions?

Plausibility ordering

Definition

m_1 is **pl-more committed** than m_2 (noted $m_1 \sqsubseteq_{pl} m_2$) if

$$Pl_1(A) \leq Pl_2(A), \quad \forall A \subseteq \Omega$$

or, equivalently,

$$Bel_1(A) \geq Bel_2(A), \quad \forall A \subseteq \Omega.$$

- Imprecise probability interpretation:

$$m_1 \sqsubseteq_{pl} m_2 \Leftrightarrow \mathcal{P}(m_1) \subseteq \mathcal{P}(m_2)$$

- Properties:

- ▶ Extension of set inclusion:

$$m_{[A]} \sqsubseteq_{pl} m_{[B]} \Leftrightarrow A \subseteq B$$

- ▶ Greatest element: vacuous mass function $m_?$

Commonality ordering

- If $m_1 = m \oplus m_2$ for some m , and if there is no conflict between m and m_2 , then $Q_1(A) = Q(A)Q_2(A) \leq Q_2(A)$ for all $A \subseteq \Omega$
- This property suggests that smaller values of the commonality function are associated with richer information content of the mass function

Definition

m_1 is **q-more committed** than m_2 (noted $m_1 \sqsubseteq_q m_2$) if

$$Q_1(A) \leq Q_2(A), \quad \forall A \subseteq \Omega$$

Properties:

- Extension of set inclusion:

$$m_{[A]} \sqsubseteq_q m_{[B]} \Leftrightarrow A \subseteq B$$

- Greatest element: vacuous mass function $m_?$

Strong (specialization) ordering

Definition

m_1 is a **specialization** of m_2 (noted $m_1 \sqsubseteq_s m_2$) if m_1 can be obtained from m_2 by distributing each mass $m_2(B)$ to subsets of B :

$$m_1(A) = \sum_{B \subseteq \Omega} S(A, B) m_2(B), \quad \forall A \subseteq \Omega,$$

where $S(A, B) =$ proportion of $m_2(B)$ transferred to $A \subseteq B$.

- S is called a **specialization matrix**
- Properties:
 - ▶ Extension of set inclusion
 - ▶ Greatest element: $m_?$
 - ▶ $m_1 \sqsubseteq_s m_2 \Rightarrow \begin{cases} m_1 \sqsubseteq_{pl} m_2 \\ m_1 \sqsubseteq_q m_2 \end{cases}$

Least Commitment Principle

Definition (Least Commitment Principle)

When several belief functions are compatible with a set of constraints, **the least informative** according to some informational ordering (if it exists) should be selected.

A very powerful method for constructing belief functions!

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Motivation

- The basic rules \oplus and \cup assume the sources of information to be **independent**, e.g.
 - ▶ experts with nonoverlapping experience/knowledge
 - ▶ nonoverlapping datasets
- What to do in case of **dependent/overlapping evidence**?
 - ▶ Describe the nature of the interaction between sources (difficult, requires a lot of information)
 - ▶ Use a combination rule that **tolerates redundancy** in the combined information
- Such rules can be derived from the LCP using **suitable informational orderings**.

Principle

- Two sources provide mass functions m_1 and m_2 , and the sources are both considered to be reliable.
- After receiving these m_1 and m_2 , the agent's state of belief should be represented by a mass function m_{12} **more committed than m_1 , and more committed than m_2 .**
- Let $\mathcal{S}_x(m)$ be the set of mass functions m' such that $m' \sqsubseteq_x m$, for some $x \in \{p, l, q, s, \dots\}$. We thus impose that

$$m_{12} \in \mathcal{S}_x(m_1) \cap \mathcal{S}_x(m_2).$$

- According to the LCP, we should select the **x -least committed element** in $\mathcal{S}_x(m_1) \cap \mathcal{S}_x(m_2)$, **if it exists.**

Need for a new ordering relation

- The above approach works for special cases.
- Example⁴: if m_1 and m_2 are consonant, then the q -least committed element in $\mathcal{S}_q(m_1) \cap \mathcal{S}_q(m_2)$ exists and it is unique: it is the consonant mass function with commonality function $Q_{12} = \min(Q_1, Q_2)$.
- In general, neither existence nor uniqueness of a solution can be guaranteed with any of the x -orderings, $x \in \{p, q, s\}$.
- We need to define a **new ordering relation**.

⁴D. Dubois and H. Prade and Ph. Smets. New Semantics for Quantitative Possibility Theory. Proc. of ECSQARU 2001, pp 410–421, Springer Verlag, 2001.

Simple mass functions

- Definition: m is **simple mass function** if it has the following form

$$m(A) = 1 - \delta(A)$$

$$m(\Omega) = \delta(A)$$

for some $A \subset \Omega$, $A \neq \emptyset$ and $\delta(A) \in (0, 1]$.

- The quantity $w(A) = -\ln \delta(A) \geq 0$ is called the **weight of evidence** for A . Mass function m is denoted by $A^{w(A)}$.
- Property:

$$A^{w_1(A)} \oplus A^{w_2(A)} = A^{w_1(A)+w_2(A)}$$

Separable mass functions

Definition (Separable mass function)

A (normalized) mass function is **separable** if it can be written as the orthogonal sum of simple mass functions:

$$m = \bigoplus_{\emptyset \neq A \subset \Omega} A^{w(A)}$$

with $w(A) \geq 0$ for all $A \subset \Omega$, $A \neq \emptyset$.

The w -ordering

Definition

Let m_1 and m_2 be two mass functions. We say that m_1 is w -more committed than m_2 (denoted by $m_1 \sqsubseteq_w m_2$) if

$$m_1 = m_2 \oplus m.$$

for some w -separable mass function m .

How to check this condition?

Weight function

- If m is separable, the corresponding weights of evidence can be recovered as

$$w(A) = \sum_{B \supseteq A} (-1)^{|B|-|A|} \ln Q(B) \quad (1)$$

for all $A \subseteq \Omega$.

- For any **nondogmatic** mass function m , (i.e., such that $m(\Omega) > 0$), we can still define “weights” from (1), but we can have $w(A) < 0$.
- Function w is called the **weight function**.
- m can be computed from w by

$$m = \bigoplus_{\emptyset \neq A \subseteq \Omega} A^{w(A)},$$

although $A^{w(A)}$ is not a proper mass function when $w(A) < 0$.

Properties of the weight function

- m is separable iff


$$w(A) \geq 0, \quad \forall A \subset \Omega, A \neq \emptyset$$

- Dempster's rule can be computed using the w -function by

$$m_1 \oplus m_2 = \bigoplus_{\emptyset \neq A \subset \Omega} A^{w_1(A)+w_2(A)}$$

- Equivalent definition of the w -ordering⁵

$$m_1 \sqsubseteq_w m_2 \Leftrightarrow w_1(A) \geq w_2(A), \quad \forall A \subset \Omega, A \neq \emptyset.$$

⁵T. Denoeux. Conjunctive and Disjunctive Combination of Belief Functions Induced by Non Distinct Bodies of Evidence. *Artificial Intelligence*, 172:234–264, 2008. 

Cautious rule

Proposition

Let m_1 and m_2 be two nondogmatic mass functions with weight functions w_1 and w_2 . The w -least committed element in $\mathcal{S}_w(m_1) \cap \mathcal{S}_w(m_2)$ exists and is unique. It is defined by:

$$m_1 \textcircled{\wedge} m_2 = \bigoplus_{\emptyset \neq A \subset \Omega} A^{\max(w_1(A), w_2(A))}$$

Operator $\textcircled{\wedge}$ is called the (normalized) cautious rule.

Computation

Cautious rule computation

m -space		w -space
m_1	\longrightarrow	w_1
m_2	\longrightarrow	w_2
$m_1 \textcircled{\wedge} m_2$	\longleftarrow	$\max(w_1, w_2)$

Remark: we often have simple mass functions in the first place, so that the w function is readily available.

Properties of the cautious rule

- Commutative, associative
- Idempotent :

$$\forall m, \quad m \circledwedge m = m$$

- Distributivity of \oplus with respect to \circledwedge

$$\forall m_1, m_2, m_3, \quad (m_1 \oplus m_2) \circledwedge (m_1 \oplus m_3) = m_1 \oplus (m_2 \circledwedge m_3)$$

The common item of evidence m_1 is not counted twice!

- No neutral element, but $m_? \circledwedge m = m$ iff m is separable.

Basic rules

The four basic rules

Sources	independent	dependent
All reliable	\oplus	$\hat{\wedge}$
At least one reliable	\odot	$\hat{\vee}$

$\hat{\vee}$ is the bold disjunctive rule

Outline

- 1 Basic notions
 - Mass functions
 - Belief and plausibility functions
 - Dempster's rule

- 2 Selected advanced topics
 - Informational orderings
 - Cautious rule
 - **Compatible frames**

Refinement and coarsening

Example

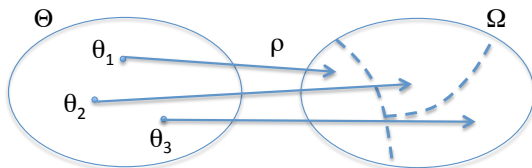
- Let us come back to the road scene analysis example, with $\Omega = \{G, R, T, O, S\}$.
- Assume that we have a **vegetation** detector, which can determine if a region of the image contains vegetation or not. For this detector, the frame of discernment is $\Theta = \{V, \neg V\}$, where V means that there is vegetation, and $\neg V$ means that there is no vegetation.
- We have the correspondence

$$\begin{aligned} V &\rightarrow \{G, T\} \\ \neg V &\rightarrow \{R, O, S\} \end{aligned}$$

- The elements of Ω can be obtained by splitting some or all of the elements of Θ . We say that Ω is a **refinement** of Θ , and Θ is a **coarsening** of Ω

Refinement and coarsening

General definition



Definition

A frame Ω is a **refinement** of a frame Θ iff there is a mapping $\rho : 2^\Theta \rightarrow 2^\Omega$ (called a **refining**) such that:

- $\{\rho(\{\theta\}), \theta \in \Theta\} \subseteq 2^\Omega$ is a partition of Ω , and
- For all $A \subseteq \Omega$, $\rho(A) = \bigcup_{\theta \in A} \rho(\{\theta\})$.

Vacuous extension

- In the road scene example, assume that the vegetation detector provides the following mass function on Θ :

$$m^\Theta(\{V\}) = 0.6, \quad m^\Theta(\{\neg V\}) = 0.3, \quad m^\Theta(\Theta) = 0.1$$

- How to express m^Θ in Ω ?
- Solution: for all $A \subseteq \Theta$, we transfer the mass $m^\Theta(A)$ to $\rho(A)$. Here,

$$\begin{aligned} m^\Theta(\{V\}) = 0.6 &\rightarrow \rho(\{V\}) = \{G, T\} \\ m^\Theta(\{\neg V\}) = 0.3 &\rightarrow \rho(\{\neg V\}) = \{R, O, S\} \\ m^\Theta(\Theta) = 0.1 &\rightarrow \rho(\Theta) = \Omega \end{aligned}$$

- We finally get the following mass function on Ω ,

$$m^{\Theta \uparrow \Omega}(\{G, T\}) = 0.6, \quad m^{\Theta \uparrow \Omega}(\{R, O, S\}) = 0.3, \quad m^{\Theta \uparrow \Omega}(\Omega) = 0.1.$$

- $m^{\Theta \uparrow \Omega}$ is called the **vacuous extension** of m^Θ in Ω .

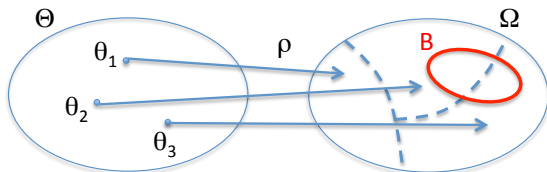
Expression of information in a coarser frame

- Let us now assume that we have the following mass function on Ω ,

$$m^\Omega(\{T\}) = 0.4, \quad m^\Omega(\{T, O\}) = 0.3, \quad m^\Omega(\{R, S\}) = 0.3.$$

- How to express m^Ω in Θ ?
- We cannot do it without loss of information, because, for instance, there is no $A \subseteq \Theta$ such that $\rho(A) = \{T\}$: the mapping ρ does not have an inverse.

Inner and outer reductions



- We can approximate any subset B of Ω by two subsets in Θ :
 - ▶ The **inner reduction** of B :

$$\underline{\rho}^{-1}(B) = \{\theta \in \Theta : \rho(\{\theta\}) \subseteq B\}$$

- ▶ The **outer reduction** of B :

$$\bar{\rho}^{-1}(B) = \{\theta \in \Theta : \rho(\{\theta\}) \cap B \neq \emptyset\}.$$

- In the example:

$$\underline{\rho}^{-1}(\{T\}) = \underline{\rho}^{-1}(\{T, O\}) = \underline{\rho}^{-1}(\{R, S\}) = \emptyset$$

$$\bar{\rho}^{-1}(\{T\}) = \{V\}, \quad \bar{\rho}^{-1}(\{T, O\}) = \{V, \neg V\}, \quad \bar{\rho}^{-1}(\{R, S\}) = \{\neg V\}$$

Restriction

Definition

The **restriction** of m^Ω in Θ is obtained by transferring each mass $m^\Omega(B)$ to the **outer reduction** of B : for all subset A of Θ ,

$$m^{\Omega \downarrow \Theta}(A) = \sum_{\bar{\rho}^{-1}(B)=A} m^\Omega(B)$$

- In the example, we thus have

$$m^{\Omega \downarrow \Theta}(\{V\}) = 0.4, \quad m^{\Omega \downarrow \Theta}(\Theta) = 0.3, \quad m^{\Omega \downarrow \Theta}(\{\neg V\}) = 0.3$$

- Remark: the vacuous extension of $m^{\Omega \downarrow \Theta}$ is

$$m^{(\Omega \downarrow \Theta) \uparrow \Omega}(\{G, T\}) = 0.4, \quad m^{(\Omega \downarrow \Theta) \uparrow \Omega}(\Omega) = 0.3$$

$$m^{(\Omega \downarrow \Theta) \uparrow \Omega}(\{R, S, O\}) = 0.3$$

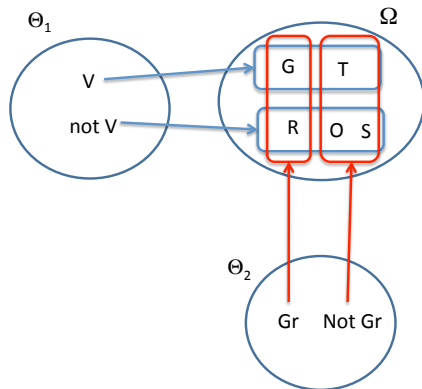
It is less precise that m^Ω : **we have lost information** when expressing m^Ω in a coarser frame.

Compatible frames of discernment

Definition

Two frames are **compatible** if they have a common refinement.

Example:



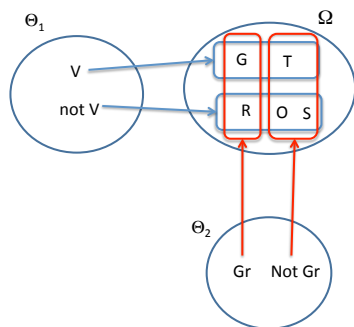
Combination of mass functions on compatible frames

Definition

Let m^{Θ_1} and m^{Θ_2} be two mass functions defined on compatible frames Θ_1 and Θ_2 with common refinement Ω . Their orthogonal sum in Ω is

$$m^{\Theta_1} \oplus m^{\Theta_2} = m^{\Theta_1 \uparrow \Omega} \oplus m^{\Theta_2 \uparrow \Omega}$$

Example



Let

$$m^{\Theta_1}(\{V\}) = 0.3, m^{\Theta_1}(\{\neg V\}) = 0.5,$$

$$m^{\Theta_1}(\{V, \neg V\}) = 0.2$$

and

$$m^{\Theta_2}(\{Gr\}) = 0.4, m^{\Theta_2}(\{\neg Gr\}) = 0.5,$$

$$m^{\Theta_2}(\{Gr, \neg Gr\}) = 0.1$$

Their extensions are

$$m^{\Theta_1 \uparrow \Omega}(\{G, T\}) = 0.3, m^{\Theta_1 \uparrow \Omega}(\{R, O, S\}) = 0.5, m^{\Theta_1 \uparrow \Omega}(\Omega) = 0.2$$

and

$$m^{\Theta_2 \uparrow \Omega}(\{G, R\}) = 0.4, m^{\Theta_2 \uparrow \Omega}(\{T, O, S\}) = 0.5, m^{\Theta_2 \uparrow \Omega}(\Omega) = 0.1$$

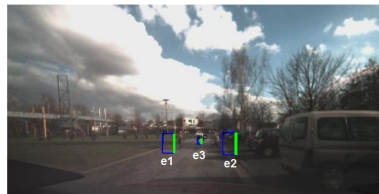
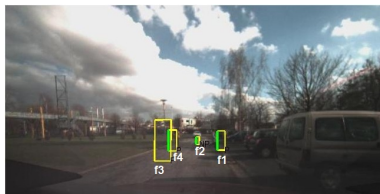
Example (continued)

Calculation of the orthogonal sum:

		$m^{\Theta_2 \uparrow \Omega}$		
		$\{G, R\}, 0.4$	$\{T, O, S\}, 0.5$	$\Omega, 0.1$
$m^{\Theta_1 \uparrow \Omega}$	$\{G, T\}, 0.3$	$\{G\}, 0.12$	$\{T\}, 0.15$	$\{G, T\}, 0.03$
	$\{R, O, S\}, 0.5$	$\{R\}, 0.2$	$\{O, S\}, 0.25$	$\{R, O, S\}, 0.05$
	$\Omega, 0.2$	$\{G, R\}, 0.08$	$\{T, O, S\}, 0.1$	$\Omega, 0.02$

Example: object association

- Let $E = \{e_1, \dots, e_n\}$ and $F = \{f_1, \dots, f_p\}$ be **two sets of objects** perceived by two sensors, or by one sensor at two different times.
- Problem: given information about each object (position, velocity, class, etc.), **find a matching between the two sets**, in such a way that each object in one set is matched with at most one object in the other set.



Method of approach

- 1 For each pair of objects $(e_i, f_j) \in E \times F$, use **sensor information** to build a pairwise mass function $m^{\Theta_{ij}}$ on the frame $\Theta_{ij} = \{h_{ij}, \bar{h}_{ij}\}$, where
 - ▶ $h_{ij} \equiv$ “ e_i and f_j are the same object”, and
 - ▶ $\bar{h}_{ij} \equiv$ “ e_i and f_j are different objects”
- 2 **Combine** the np mass functions $m^{\Theta_{ij}}$
- 3 Find the matching relation with the **highest plausibility**.

Building the pairwise mass functions

Using position information

- Assume that each sensor provides an **estimated position** for each object. Let d_{ij} denote the distance between the estimated positions of e_i and f_j , computed using some distance measure.
- A small value of d_{ij} supports hypothesis h_{ij} , while a large value of d_{ij} supports hypothesis \bar{h}_{ij} . Depending on sensor reliability, a fraction of the unit mass should also be assigned to $\Theta_{ij} = \{h_{ij}, \bar{h}_{ij}\}$.
- Model:

$$m_p^{\Theta_{ij}}(\{h_{ij}\}) = \alpha \varphi(d_{ij})$$

$$m_p^{\Theta_{ij}}(\{\bar{h}_{ij}\}) = \alpha (1 - \varphi(d_{ij}))$$

$$m_p^{\Theta_{ij}}(\Theta_{ij}) = 1 - \alpha$$

where $\alpha \in [0, 1]$ is a degree of confidence in the sensor information and φ is a decreasing function taking values in $[0, 1]$.

Building the pairwise mass functions

Using velocity information

- Let us now assume that each sensor returns a **velocity vector** for each object. Let d'_{ij} denote the distance between the velocities of objects e_i and f_j .
- Here, a large value of d'_{ij} supports the hypothesis \bar{h}_{ij} , whereas a small value of d'_{ij} does not support specifically h_{ij} or \bar{h}_{ij} , as two distinct objects may have similar velocities.
- Model:

$$m_v^{\Theta_{ij}}(\{\bar{h}_{ij}\}) = \alpha' \psi(d'_{ij})$$

$$m_v^{\Theta_{ij}}(\Theta_{ij}) = 1 - \alpha' \psi(d'_{ij}),$$

where $\alpha' \in [0, 1]$ is a degree of confidence in the sensor information and ψ is an increasing function taking values in $[0, 1]$.

Building the pairwise mass functions

Using class information

- Let us assume that the **objects belong to classes**. Let Ω be the set of possible classes, and let m_i and m_j denote mass functions representing evidence about the class membership of objects e_i and f_j .
- If e_i and f_j do not belong to the same class, they cannot be the same object. However, if e_i and f_j do belong to the same class, they may or may not be the same object.
- We can show that the mass function $m_c^{\Theta_{ij}}$ on Θ_{ij} derived from m_i and m_j has the following expression:

$$m_c^{\Theta_{ij}}(\{\bar{h}_{ij}\}) = \kappa_{ij}$$
$$m_c^{\Theta_{ij}}(\Theta_{ij}) = 1 - \kappa_{ij},$$

where κ_{ij} is the **degree of conflict** between m_i and m_j

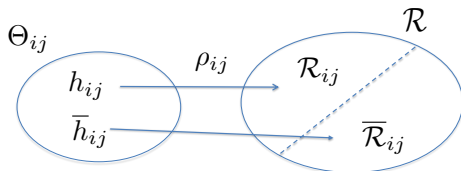
Combination

- For each object pair (e_i, f_j) , a **pairwise mass function** $m^{\Theta_{ij}}$ representing all the available evidence about Θ_{ij} can finally be obtained as:

$$m^{\Theta_{ij}} = m_p^{\Theta_{ij}} \oplus m_v^{\Theta_{ij}} \oplus m_c^{\Theta_{ij}}$$

- How to combine the np mass functions $m^{\Theta_{ij}}$?
- Does there exist a common refinement of the frames Θ_{ij} for (i, j) ?

Common refinement



- Let \mathcal{R} be the set of all admissible matching relations, and let $\mathcal{R}_{ij} \subseteq \mathcal{R}$ be the subset of relations R such that $(e_i, f_j) \in R$.
- We can define a refining ρ_{ij} from $2^{\Theta_{ij}}$ to $2^{\mathcal{R}}$. The frames Θ_{ij} are compatible.
- Vacuously extending $m^{\Theta_{ij}}$ in \mathcal{R} yields the following mass function:

$$m^{\Theta_{ij} \uparrow \mathcal{R}}(\mathcal{R}_{ij}) = m^{\Theta_{ij}}(\{h_{ij}\}) = \alpha_{ij}$$

$$m^{\Theta_{ij} \uparrow \mathcal{R}}(\bar{\mathcal{R}}_{ij}) = m^{\Theta_{ij}}(\{\bar{h}_{ij}\}) = \beta_{ij}$$

$$m^{\Theta_{ij} \uparrow \mathcal{R}}(\mathcal{R}) = m^{\Theta_{ij}}(\Theta_{ij}) = 1 - \alpha_{ij} - \beta_{ij}.$$

Combination of contour functions

- The frame \mathcal{R} is very big and computing the orthogonal sum of the np mass functions

$$m^{\mathcal{R}} = \bigoplus_{i,j} m^{\Theta_{ij} \uparrow \mathcal{R}}$$

has exponential complexity.

- Instead, we will only compute the combined contour function ρ_l corresponding to $m^{\mathcal{R}}$. We recall that

$$\rho_l \propto \prod_{i,j} \rho_{lij},$$

where ρ_{lij} denote the **contour function** corresponding to $m^{\Theta_{ij} \uparrow \mathcal{R}}$.

Expression of contour functions

- We have

$$m^{\Theta_{ij} \uparrow \mathcal{R}}(\mathcal{R}_{ij}) = \alpha_{ij}, \quad m^{\Theta_{ij} \uparrow \mathcal{R}}(\overline{\mathcal{R}_{ij}}) = \beta_{ij}, \quad m^{\Theta_{ij} \uparrow \mathcal{R}}(\mathcal{R}) = 1 - \alpha_{ij} - \beta_{ij}.$$

- For all $R \in \mathcal{R}$,

$$\begin{aligned} pl_{ij}(R) &= \begin{cases} 1 - \beta_{ij} & \text{if } R \in \mathcal{R}_{ij}, \\ 1 - \alpha_{ij} & \text{otherwise,} \end{cases} \\ &= (1 - \beta_{ij})^{R_{ij}} (1 - \alpha_{ij})^{1 - R_{ij}}, \end{aligned}$$

where $R_{ij} = 1$ if e_i and f_j are matched and $R_{ij} = 0$ otherwise.

- Consequently, the combined contour function is

$$pl(R) \propto \prod_{i,j} (1 - \beta_{ij})^{R_{ij}} (1 - \alpha_{ij})^{1 - R_{ij}}.$$

Finding the most plausible matching

- We have

$$\begin{aligned} \ln pl(R) &= \sum_{i,j} [R_{ij} \ln(1 - \beta_{ij}) + (1 - R_{ij}) \ln(1 - \alpha_{ij})] + C \\ &= \sum_{i,j} R_{ij} \ln \frac{1 - \beta_{ij}}{1 - \alpha_{ij}} + C' \end{aligned}$$

- The **most plausible relation** R^* can thus be found by solving the following **binary linear optimization problem**:

$$\max \sum_{i,j} R_{ij} \ln \frac{1 - \beta_{ij}}{1 - \alpha_{ij}}$$

subject to $R_{ij} \in \{0, 1\}$, $\forall(i, j)$, $\sum_{j=1}^p R_{ij} \leq 1$, $\forall i$ and $\sum_{i=1}^n R_{ij} \leq 1$, $\forall j$.

- This problem can be shown to be equivalent to a **linear assignment problem** and can be solved in $O(\max(n, p)^3)$ time.

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cf. <http://www.hds.utc.fr/~tdenoeux>



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